Chapter Two: Linear Programming: Model Formulation and Graphical Solution

PROBLEM SUMMARY

- 1. Maximization
- 2. Maximization
- 3. Minimization
- **4.** Sensitivity analysis (2–3)
- 5. Minimization
- 6. Maximization
- 7. Slack analysis (2–6)
- 8. Sensitivity analysis (2–6)
- 9. Maximization, graphical solution
- **10.** Slack analysis (2–9)
- 11. Maximization, graphical solution
- **12.** Minimization, graphical solution
- 13. Maximization, graphical solution
- **14.** Sensitivity analysis (2–13)
- **15.** Sensitivity analysis (2–13)
- 16. Maximization, graphical solution
- **17.** Sensitivity analysis (2–16)
- 18. Maximization, graphical solution
- 19. Standard form
- 20. Maximization, graphical solution
- 21. Standard form
- 22. Maximization, graphical solution
- **23.** Constraint analysis (2–22)
- 24. Minimization, graphical solution
- **25.** Sensitivity analysis (2–24)
- **26.** Sensitivity analysis (2–24)
- **27.** Sensitivity analysis (2–24)
- **28.** Minimization, graphical solution
- **29.** Minimization, graphical solution
- **30.** Sensitivity analysis (2–29)
- 31. Maximization, graphical solution
- 32. Maximization, graphical solution
- 33. Minimization, graphical solution
- 34. Maximization, graphical solution

- **35.** Sensitivity analysis (2–34)
- **36.** Minimization, graphical solution
- **37.** Maximization, graphical solution
- **38.** Maximization, graphical solution
- **39.** Sensitivity analysis (2–38)
- **40.** Maximization, graphical solution
- **41.** Sensitivity analysis (2–40)
- **42.** Maximization, graphical solution
- **43.** Sensitivity analysis (2–42)
- 44. Minimization, graphical solution
- **45.** Sensitivity analysis (2–44)
- **46.** Maximization, graphical solution
- **47.** Sensitivity analysis (2–46)
- **48.** Maximization, graphical solution
- **49.** Sensitivity analysis (2–48)
- **50.** Multiple optimal solutions
- **51.** Infeasible problem
- **52.** Unbounded problem

PROBLEM SOLUTIONS



2.(a) maximize $Z = 6x_1 + 4x_2$ (profit, \$) subject to

 $10x_1 + 10x_2 \le 100$ (line 1, hr) $7x_1 + 3x_2 \le 42$ (line 2, hr) $x_1, x_2 \ge 0$

Wood

$$2x_1 + 6x_2 \le 36 \text{ lb}$$

$$2(6) + 6(3.2) \le 36$$

$$12 + 19.2 \le 36$$

$$31.2 \le 36$$

$$36 - 31.2 = 4.8$$

There is 4.8 lb of wood left unused.

- 8. The new objective function, $Z = 400x_1 + 500x_2$, is parallel to the constraint for labor, which results in multiple optimal solutions. Points *B* $(x_1 = 30/7, x_2 = 32/7)$ and *C* $(x_1 = 6, x_2 = 3.2)$ are the alternate optimal solutions, each with a profit of \$4,000.
- **9.(a)** maximize $Z = x_1 + 5x_2$ (profit, \$) subject to

 $5x_1 + 5x_2 \le 25 \text{ (flour, lb)}$ $2x_1 + 4x_2 \le 16 \text{ (sugar, lb)}$ $x_1 \le 5 \text{ (demand for cakes)}$ $x_1, x_2 \ge 0$



10. In order to solve this problem, you must substitute the optimal solution into the resource constraints for flour and sugar and determine how much of each resource is left over.

Flour

$$5x_1 + 5x_2 \le 25$$
 lb
 $5(0) + 5(4) \le 25$
 $20 \le 25$
 $25 - 20 = 5$

There are 5 lb of flour left unused.

Sugar

 $\begin{array}{l} 2x_1 + 4x_2 \leq 16 \\ 2(0) + 4(4) \leq 16 \\ 16 \leq 16 \end{array}$

There is no sugar left unused.





$$3x_1 + x_2 \ge 6 \text{ (antibiotic 1, units)}$$

$$x_1 + x_2 \ge 4 \text{ (antibiotic 2, units)}$$

$$2x_1 + 6x_2 \ge 12 \text{ (antibiotic 3, units)}$$

$$x_1, x_2 \ge 0$$



13.(a) maximize $Z = 300x_1 + 400x_2$ (profit, \$) subject to

$$3x_1 + 2x_2 \le 18 \text{ (gold, oz)}$$

$$2x_1 + 4x_2 \le 20 \text{ (platinum, oz)}$$

$$x_2 \le 4 \text{ (demand, bracelets)}$$

$$x_1, x_2 \ge 0$$

The extreme points to evaluate are now A, B', and C'.

A:
$$x_1 = 0$$

 $x_2 = 30$
 $Z = 1,200$
*B': $x_1 = 15.8$
 $x_2 = 20.5$
 $Z = 1,610$
C': $x_1 = 24$
 $x_2 = 0$

$$Z = 1,200$$

Point B' is optimal





 $x_1 + s_2 = 4$ $x_2 + s_2 = 6$ $x_1 + x_2 + s_3 = 5$ $x_1, x_2 \ge 0$ A: $s_1 = 4, s_2 = 1, s_3 = 0$ B: $s_1 = 0, s_2 = 5, s_3 = 0$ C: $s_1 = 0, s_2 = 6, s_3 = 1$





21. maximize $Z = 5x_1 + 8x_2 + 0s_1 + 0s_3 + 0s_4$ subject to:

> $3x_1 + 5x_2 + s_1 = 50$ $2x_1 + 4x_2 + s_2 = 40$ $x_1 + s_3 = 8$ $x_1, x_2 \ge 0$ $a_1 = 0, a_2 = 8, a_3 = 0$

A:
$$s_1 = 0$$
, $s_2 = 0$, $s_3 = 8$, $s_4 = 0$
B: $s_1 = 0$, $s_2 = 3.2$, $s_3 = 0$, $s_4 = 4.8$
C: $s_1 = 26$, $s_2 = 24$, $s_3 = 0$, $s_4 = 10$

22.



23. It changes the optimal solution to point *A* $(x_1 = 8, x_2 = 6, Z = 112)$, and the constraint, $x_1 + x_2 \le 15$, is no longer part of the solution space boundary.

24.(a) Minimize $Z = 64x_1 + 42x_2$ (labor cost, \$) subject to $16x_1 + 12x_2 \ge 450$ (claims) $x_1 + x_2 \le 40$ (workstations) $0.5x_1 + 1.4x_2 \le 25$ (defective claims) $x_1, x_2 \ge 0$



- **25.** Changing the pay for a full-time claims processor from \$64 to \$54 will change the solution to point A in the graphical solution where $x_1 = 28.125$ and $x_2 = 0$, i.e., there will be no part-time operators. Changing the pay for a part-time operator from \$42 to \$36 has no effect on the number of full-time and part-time operators hired, although the total cost will be reduced to \$1,671.95
- 26. Eliminating the constraint for defective claims would result in a new solution, $x_1 = 0$ and $x_2 = 37.5$, where only part-time operators would be hired.
- **27.** The solution becomes infeasible; there are not enough workstations to handle the increase in the volume of claims.





30. The problem becomes infeasible.



47. A new constraint is added to the model in

$$\frac{x_1}{x_2} \ge 1.5$$

The solution is, $x_1 = 160$, $x_2 = 106.67$, Z = \$568





 $\begin{aligned} x_1 + x_2 &\leq 50 \text{ (available land, acres)} \\ 10x_1 + 3x_2 &\leq 300 \text{ (labor, hr)} \\ 8x_1 + 20x_2 &\leq 800 \text{ (fertilizer, tons)} \\ x_1 &\leq 26 \text{ (shipping space, acres)} \\ x_2 &\leq 37 \text{ (shipping space, acres)} \\ x_1, x_2 &\geq 0 \end{aligned}$



49. The feasible solution space changes if the fertilizer constraint changes to $20x_1 + 20x_2 \le 800$ tons. The new solution space is A'B'C'D'. Two of the constraints now have no effect.



The new optimal solution is point C':





Multiple optimal solutions; A and B alternate optimal



CASE SOLUTION: METROPOLITAN POLICE PATROL

The linear programming model for this case problem is minimize Z = x/60 + y/45

subject to

$$2x + 2y \ge 5$$
$$2x + 2y \le 12$$
$$y \ge 1.5x$$
$$x, y \ge 0$$

The objective function coefficients are determined by dividing the distance traveled, ie., x/3, by the travel speed, ie., 20 mph. Thus, the *x* coefficient is $x/3 \div 20$, or x/60. In the first two constraints, 2x + 2y represents the formula for the perimeter of a rectangle.

The graphical solution is displayed as follows.



The optimal solution is x = 1, y = 1.5, and Z = 0.05. This means that a patrol sector is 1.5 miles by 1 mile and the response time is 0.05 hr, or 3 min.

CASE SOLUTION: "THE POSSIBILITY" RESTAURANT

The linear programming model formulation is

Maximize = $Z = \$12x_1 + 16x_2$

subject to

$$x_1 + x_2 \le 60$$

$$.25x_1 + .50x_2 \le 20$$

$$x_1/x_2 \ge 3/2 \text{ or } 2x_1 - 3x_2 \ge 0$$

$$x_2/(x_1 + x_2) \le .10 \text{ or } .90x_2 - .10x_1 \ge 0$$

$$x_1, x_2 \ge 0$$

The graphical solution is shown as follows.

